

# CS101 Midterm Exam

Fall 2007

October 15, 2007

## I. Search

### 1. Terms and Definitions.

Define in your own words the following terms associated with search:

1. state:
2. state space:
3. search tree:
4. search node:
5. goal:
6. action:
7. successor function:
8. branching factor:

### 2. Search.

Suppose there are two friends living in different cities on a map, such as the Romania map in the textbook. On every turn, we can move each friend simultaneously to a neighboring city on the map. The amount of time needed to move from city  $i$  to neighbor  $j$  is equal to the road distance  $d(i, j)$  between the cities, but on each turn the friend that arrives first must wait until the other one arrives (and calls the first on his/her cell phone) before the next turn can begin. We want the two friends to meet as quickly as possible. Let us formulate this as a search problem.

- (a) What is the state space? (You will find it helpful to define some formal notation here.)
- (b) What is the successor function?
- (c) What is the goal?
- (d) Let  $SLD(i, j)$  be the straight-line distance between any two cities  $i$  and  $j$ . Which, if any, of the following heuristic functions are admissible? (If none, write NONE.)
  - (i)  $SLD(i, j)$
  - (ii)  $2 \times SLD(i, j)$
  - (iii)  $SLD(i, j)/2$
- (e) True/False: There are completely connected maps for which no solution exists.

#### 4. Constraint Satisfaction Problems

A constraint satisfaction problem (or CSP) is defined by a set of variables,  $x_1, x_2, \dots, x_n$ , and a set of constraints,  $c_1, c_2, \dots, c_m$ . Each variable  $x_i$  has a nonempty domain  $D_i$  of possible values. Each constraint  $c_i$  involves some subset of the variables and specifies the allowable combinations of values for that subset. A state of the problem is defined by an assignment of values to some or all of the variables. Now, give precise formulations for each as constraint satisfaction problems:

1. Rectilinear floor planning: find non-overlapping places in a large rectangle for a number of smaller rectangles.
2. Class scheduling: There is a fixed number of professors and classrooms, a list of classes to be offered, and a list of possible time slots for classes. Each professor has a set of classes that he or she can teach.

## II. Logic and Reasoning

### 1. Natural Deduction

1. Prove  $(p \wedge q) \wedge r, s \wedge t \vdash q \wedge s$
2. Prove  $(p \rightarrow r) \wedge (q \rightarrow r) \vdash (p \wedge q) \rightarrow r$

### 2. Translations from English to First-order logic

1. No man is both a butcher and a vegetarian.
2. All men except butchers like vegetarians.
3. The only vegetarian butchers are women.
4. No man likes a woman who is a vegetarian.

### 3. Translations to Conjunctive Clausal Form

A formula is in conjunctive normal form (CNF) if it is a conjunction of clauses, where a clause is a disjunction of literals. For example, all of the following formulas are in CNF:  $A \wedge B$ ,  $\neg A \wedge (B \vee C)$ , and  $(A \vee B) \wedge (\neg B \vee C)$ .

1.  $\forall x \forall y [P(x, y) \rightarrow [Q(x, y) \rightarrow R(x, y)]]$
2.  $\neg \forall x \exists y [P(x, y) \rightarrow Q(x, y)]$
3.  $\forall x \forall y [(P(x, y) \vee Q(x, y)) \rightarrow R(x, y)]$

### 4. Unification

For each pair, give the substitution that unifies the two expressions:

1.  $\text{Color}(\text{Tweety}, \text{Yellow}), \text{Color}(x, y)$
2.  $\text{Color}(\text{Hat}(\text{Mailman}), \text{Blue}), \text{Color}(\text{Hat}(y), x)$
3.  $R(\text{F}(x), B), R(y, z)$
4.  $R(\text{F}(y), x), R(x, \text{F}(B))$

### Extra Credit 1.

Solve by resolution to show that I win.

Heads I win, tails you lose.

Here are the axioms and expression to prove, where  $H$  is "heads",  $W$  is "I win",  $T$  is "tails", and  $L$  is "You lose".

1.  $H \rightarrow W, T \rightarrow L, T \vee H, L \rightarrow W \vdash W$

### Extra Credit 2.

Translate the following from First-order logic to English.

1.  $\forall x[\text{hesitates}(x) \rightarrow \text{lost}(x)]$
2.  $\neg \exists x[\text{business}(x) \wedge \text{is-like}(x, \text{Showbusiness})]$
3.  $\exists x \forall t[\text{Person}(x) \wedge \text{time}(t) \wedge \text{Can-fool}(x, t)]$

## Appendix: Strategies for Natural Deductive Proofs

Given a set of premises,  $\Delta$ , and the goal you are trying to prove,  $\Gamma$ , there are some simple rules of thumb that will be helpful in finding ND proofs for problems. These are not complete, but will get you through your problems. Here goes:

Given  $\Delta$ , prove  $\Gamma$ :

1. Apply premises,  $p_i \in \Delta$ , to prove  $\Gamma$ .
2. If you need to use a  $\vee$ -premise, then apply  $\vee$ -e and prove  $\Gamma$  for each disjunct.
3. Otherwise, work backwards from the type of goal you are proving:
  - (a) If the goal  $\Gamma$  is a conditional ( $A \rightarrow B$ ), then assume  $A$  and prove  $B$ ; use  $\rightarrow$ -i rule.
  - (b) If the goal  $\Gamma$  is a negative ( $\neg A$ ), then assume  $\neg \neg A$  and prove contradiction; use  $\neg$ -i rule.
  - (c) If the goal  $\Gamma$  is a conjunction ( $A \wedge B$ ), then prove  $A$  and prove  $B$ ; use  $\wedge$ -i rule.
  - (d) If the goal  $\Gamma$  is a disjunction ( $A \vee B$ ), then prove one of  $A$  or  $B$ ; use  $\vee$ -i rule.